

# CTY LIN: Markov Chains and the Google Ranking System

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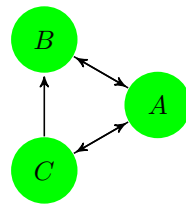
## 1 Matrices of Google

*Reminders:* In 1996, Google founders Sergey Brin and Larry Page developed the PageRank algorithm, based on the amount of links and page has going to and from it. A page with many inbound links is ranked higher, and a page that is off in the middle of nowhere will be ranked lower.

*Note:* The notation for the page rank,  $r(A)$ , will be shortened to just  $A$  in this paper.

### 1.1 Finding PageRanks

We have the websites  $A$ ,  $B$ , and  $C$  and their web of connections.



We see that

$$\begin{aligned} A &= B + \frac{1}{2}C \\ B &= \frac{1}{2}A + \frac{1}{2}C \\ C &= \frac{1}{2}A \end{aligned}$$

$$\begin{bmatrix} 0 & 1 & 0.5 \\ 0.5 & 0 & 0.5 \\ 0.5 & 0 & 0 \end{bmatrix} \vec{r} = \vec{r}$$

$$\begin{bmatrix} -1 & 1 & 0.5 \\ 0.5 & -1 & 0.5 \\ 0.5 & 0 & -1 \end{bmatrix} \vec{r} = \vec{0}$$

We now row reduce this matrix:

$$\begin{bmatrix} -1 & 1 & 0.5 \\ 0.5 & -1 & 0.5 \\ 0.5 & 0 & -1 \end{bmatrix} \text{ Flip } R_3 \text{ and } R_1 = \begin{bmatrix} 0.5 & 0 & -1 \\ 0.5 & -1 & 0.5 \\ -1 & 1 & 0.5 \end{bmatrix} R_3 = R_3 + 2R_1, R_2 = R_2 - R_1$$

$$= \begin{bmatrix} 0.5 & 0 & -1 \\ 0 & -1 & 1.5 \\ 0 & 1 & -1.5 \end{bmatrix} R_3 = R_3 + R_2$$

$$= \begin{bmatrix} 0.5 & 0 & -1 \\ 0 & -1 & 1.5 \\ 0 & 0 & 0 \end{bmatrix}$$

This matrix tells us that  $\frac{1}{2}A = C$  and  $B = \frac{3}{2}C$ . If  $C = 2$ , then  $B = 3$  and  $A = 4$ , giving us the vector

$$\begin{bmatrix} 4 \\ 3 \\ 2 \end{bmatrix}$$

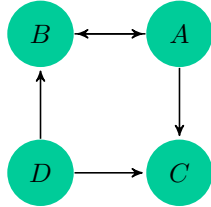
But because the entries must add to one, we divide by  $2 + 3 + 4 = 9$  to get

$$\begin{bmatrix} \frac{4}{9} \\ \frac{3}{9} \\ \frac{2}{9} \end{bmatrix}$$

The order of the pages on Google will be  $A, B, C$ .

## 1.2 Altered Transition Matrices

Sometimes, however, these matrices don't always work. If a webpage only has inbound links, then we'll get trapped there eventually, as is in the case of the following diagram.



$$A = B$$

$$B = \frac{1}{2}A + \frac{1}{2}D$$

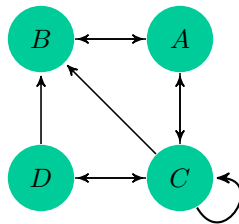
$$C = \frac{1}{2}A + \frac{1}{2}D$$

$$D = 0$$

Which has the matrix

$$\begin{bmatrix} 0 & 1 & 0 & 0 \\ 0.5 & 0 & 0 & 0.5 \\ 0.5 & 0 & 0 & 0.5 \\ 0 & 0 & 0 & 0 \end{bmatrix} \vec{r} = \vec{r}$$

The column of zeroes is reflective of the problem of being trapped on webpage  $C$ . We can change that by changing  $C$  to have outbound links to all pages including itself.



We can now build an *altered transition matrix*, based on the altered diagram

$$A = B + \frac{1}{4}C$$

$$B = \frac{1}{2}A + \frac{1}{2}D + \frac{1}{4}C$$

$$C = \frac{1}{2}A + \frac{1}{2}D + \frac{1}{4}C$$

$$D = \frac{1}{4}C$$

which has matrix

$$\begin{bmatrix} 0 & 1 & 0.25 & 0 \\ 0.5 & 0 & 0.25 & 0.5 \\ 0.5 & 0 & 0.25 & 0.5 \\ 0 & 0 & 0.25 & 0 \end{bmatrix} \vec{r} = \vec{r}$$

$$\begin{bmatrix} -1 & 1 & 0.25 & 0 \\ 0.5 & -1 & 0.25 & 0.5 \\ 0.5 & 0 & -0.75 & 0.5 \\ 0 & 0 & 0.25 & -1 \end{bmatrix} \vec{r} = 0$$

An now we row reduce and solve for  $\vec{r}$ .

$$\begin{bmatrix} -1 & 1 & 0.25 & 0 \\ 0.5 & -1 & 0.25 & 0.5 \\ 0.5 & 0 & -0.75 & 0.5 \\ 0 & 0 & 0.25 & -1 \end{bmatrix} R_2 = R_1 + 2R_2, R_3 = R_1 + 2R_3$$

$$= \begin{bmatrix} -1 & 1 & 0.25 & 0 \\ 0 & -1 & 0.75 & 1 \\ 0 & 1 & -1.25 & 1 \\ 0 & 0 & 0.25 & -1 \end{bmatrix} R_3 = R_2 + R_3$$

$$= \begin{bmatrix} -1 & 1 & 0.25 & 0 \\ 0 & -1 & 0.75 & 1 \\ 0 & 0 & -0.5 & 2 \\ 0 & 0 & 0.25 & -1 \end{bmatrix} R_4 = R_3 + 2R_4$$

$$= \begin{bmatrix} -1 & 1 & 0.25 & 0 \\ 0 & -1 & 0.75 & 1 \\ 0 & 0 & -0.5 & 2 \\ 0 & 0 & 0 & 0 \end{bmatrix} R_2 = 3R_3 + 2R_2, R_1 = 2R_1 + R_3$$

$$= \begin{bmatrix} -2 & 2 & 0 & 2 \\ 0 & -2 & 0 & 8 \\ 0 & 0 & -0.5 & 2 \\ 0 & 0 & 0 & 0 \end{bmatrix} R_2 = \frac{1}{2}R_2, R_1 = \frac{1}{2}(R_1 + R_2), R_3 = 2R_3$$

$$= \begin{bmatrix} -1 & 0 & 0 & 5 \\ 0 & -1 & 0 & 4 \\ 0 & 0 & -1 & 4 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

We can write the solution space as

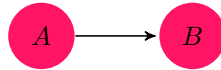
$$\begin{bmatrix} A \\ B \\ C \\ D \end{bmatrix} = \begin{bmatrix} 5D \\ 4D \\ 4D \\ D \end{bmatrix} = D \begin{bmatrix} 5 \\ 4 \\ 4 \\ 1 \end{bmatrix}$$

We can set  $D$  to be the reciprocal of  $5 + 4 + 4 + 1 = 14$  to get the PageRank vector

$$\begin{bmatrix} \frac{5}{14} \\ \frac{4}{14} \\ \frac{4}{14} \\ \frac{1}{14} \end{bmatrix}$$

### 1.3 Ralph's Link Diagram

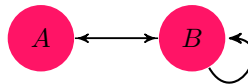
Given the link diagram,



and that Ralph clicks on a random link every minute, after four minutes, he's going to be on  $B$ .

After all, no matter the starting probability vector (in this case  $\begin{bmatrix} \frac{2}{3} \\ \frac{1}{3} \end{bmatrix}$ ), once Ralph gets to  $B$  (which he will get to within 1 minute), he is trapped.

Now, given the more interesting problem, that Ralph will have a 50% chance of going to either  $A$  or  $B$  once he lands on  $B$ , we realize this is just a clever set-up for the altered link diagram,



$$A = \frac{1}{2}B$$

$$B = A + \frac{1}{2}B$$

We aren't trying to find the PageRank vector this time, though. We are trying to find the probability vector, which means we have to interpret this as a *probability* matrix  $\tilde{P} =$

$$\begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix}$$

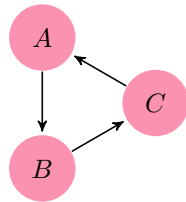
Now, after four minutes, the probability vector  $\vec{v}$  will just be the original probability vector  $\begin{bmatrix} 2 \\ 31 \\ 3 \end{bmatrix}$ , multiplied by  $\tilde{P}$  four times.

$$\begin{aligned} \vec{v} &= \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} \frac{2}{3} \\ \frac{1}{3} \end{bmatrix} \\ &= \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} \frac{1}{6} \\ \frac{5}{6} \end{bmatrix} \\ &= \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} \frac{5}{12} \\ \frac{7}{12} \end{bmatrix} \\ &= \begin{bmatrix} 0 & 0.5 \\ 1 & 0.5 \end{bmatrix} \begin{bmatrix} \frac{7}{24} \\ \frac{17}{24} \end{bmatrix} \\ &= \begin{bmatrix} \frac{17}{48} \\ \frac{31}{48} \end{bmatrix} \end{aligned}$$

The probability that Ralph will be on page  $B$  after 4 minutes is  $\frac{31}{48}$ .

## 1.4 Circle of Links

If we have the link diagram



$$A = C$$

$$B = A$$

$$C = B$$

$$\tilde{P} = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$$

Allow us to define  $\vec{x} = \begin{bmatrix} 0.2 \\ 0.3 \\ 0.5 \end{bmatrix}$ . What happens when we multiply  $\vec{x}$  with  $\tilde{P}$ ?

$$\tilde{P}\vec{x} = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0.2 \\ 0.3 \\ 0.5 \end{bmatrix} = \begin{bmatrix} 0.5 \\ 0.2 \\ 0.3 \end{bmatrix}$$

$$\tilde{P}^2\vec{x} = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0.5 \\ 0.2 \\ 0.3 \end{bmatrix} = \begin{bmatrix} 0.3 \\ 0.5 \\ 0.2 \end{bmatrix}$$

$$\tilde{P}^3\vec{x} = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0.3 \\ 0.5 \\ 0.2 \end{bmatrix} = \begin{bmatrix} 0.2 \\ 0.3 \\ 0.5 \end{bmatrix}$$

We can see that the matrix  $\tilde{P}$  just creates cycles between 3 states. Thus, ALL  $\tilde{P}^n\vec{x}$  for  $n \equiv 0 \pmod{3}$

will be  $\begin{bmatrix} 0.2 \\ 0.3 \\ 0.5 \end{bmatrix}$ , and all  $\tilde{P}^n\vec{x}$  for  $n \equiv 1 \pmod{3}$  will be  $\begin{bmatrix} 0.5 \\ 0.2 \\ 0.3 \end{bmatrix}$ , and so on. Because  $100 \equiv 1 \pmod{3}$ ,

$$\tilde{P}^{100}\vec{x} = \begin{bmatrix} 0.5 \\ 0.2 \\ 0.3 \end{bmatrix}.$$

$\tilde{P}^k\vec{x}$  will never approach  $\vec{r} = \begin{bmatrix} \frac{1}{3} \\ \frac{1}{3} \\ \frac{1}{3} \end{bmatrix}$ , because it will only oscillate, never approaching any one particular value. Also,  $\tilde{P}$  is not regular. There will always be 0's, because all it does is shuffle the rows, and the rows all have zeroes.

## 1.5 The Google Matrix

*Reminders:*  $G_\alpha = \alpha P + (1 - \alpha)M$ , where  $P$  is the transition matrix, and  $M$  is a matrix with each entry as  $\frac{1}{n}$  where  $n$  is the dimension of the matrix.

For the matrix above, we have  $G_{0.5}$  as

$$\frac{1}{2} \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} + \frac{1}{2} \begin{bmatrix} \frac{1}{3} & \frac{1}{3} & \frac{1}{3} \\ \frac{1}{3} & \frac{1}{3} & \frac{1}{3} \\ \frac{1}{3} & \frac{1}{3} & \frac{1}{3} \end{bmatrix} = \begin{bmatrix} \frac{1}{6} & \frac{1}{6} & \frac{2}{3} \\ \frac{2}{3} & \frac{1}{6} & \frac{1}{6} \\ \frac{1}{6} & \frac{2}{3} & \frac{1}{6} \end{bmatrix}$$

The vector  $\begin{bmatrix} \frac{1}{3} \\ \frac{1}{3} \\ \frac{1}{3} \end{bmatrix}$  ( $\vec{r}$ ) is also a steady-state vector of this Google matrix.

$$\begin{bmatrix} \frac{1}{6} & \frac{1}{6} & \frac{2}{3} \\ \frac{2}{3} & \frac{1}{6} & \frac{1}{6} \\ \frac{1}{6} & \frac{2}{3} & \frac{1}{6} \end{bmatrix} \begin{bmatrix} \frac{1}{3} \\ \frac{1}{3} \\ \frac{1}{3} \end{bmatrix} = \begin{bmatrix} \frac{1}{18} + \frac{1}{18} + \frac{4}{18} \\ \frac{4}{18} + \frac{1}{18} + \frac{1}{18} \\ \frac{1}{18} + \frac{4}{18} + \frac{1}{18} \end{bmatrix} = \begin{bmatrix} \frac{1}{3} \\ \frac{1}{3} \\ \frac{1}{3} \end{bmatrix}$$

Also,  $G_{0.5}^5 \vec{x}$  (same  $\vec{x}$  as above), is very close to  $\vec{r}$ . Because the Google matrix is a regular stochastic matrix,  $G_{\alpha}^k \vec{x}$  converges to the steady-state vector of  $G_{\alpha}$ , which in this case is  $\vec{r}$ .